

Optimizing the Barnes-Hut Algorithm for Multicore Clusters

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Motivation

- Important trends of supercomputer evolution:
 - Number of cores per node keeps increasing;
 - Amount of memory per core is decreasing;
 - One-sided communication (rDMA) is increasingly well supported on the interconnection networks.
- It is important to explore new programming models to cope with or leverage these trends.
- Through the Barnes-Hut (BH) algorithm, we study a promising programming paradigm, i.e., partitioned global address space (PGAS) + X, which
 - Uses a global address space for internode communication;
 - Hides latency through intranode multitasking;
 - Achieves data sharing and message reduction through multithreading.
- We show many advantages gained from this design.

What is Barnes-Hut?

- An $O(n \log n)$ algorithm for the n-body problem, which needs to compute forces on every body (particle).
- Use an octree as the main data structure.
- Has **massive** parallelisms:
 - For all bodies, can compute their forces in parallel.
 - For one body, can interact with different tree nodes (aka *cells*) in parallel, provided the results are summed together finally.
- Irregular**, whether to traverse children of a cell (i.e., *open* a cell) is computation dependent.

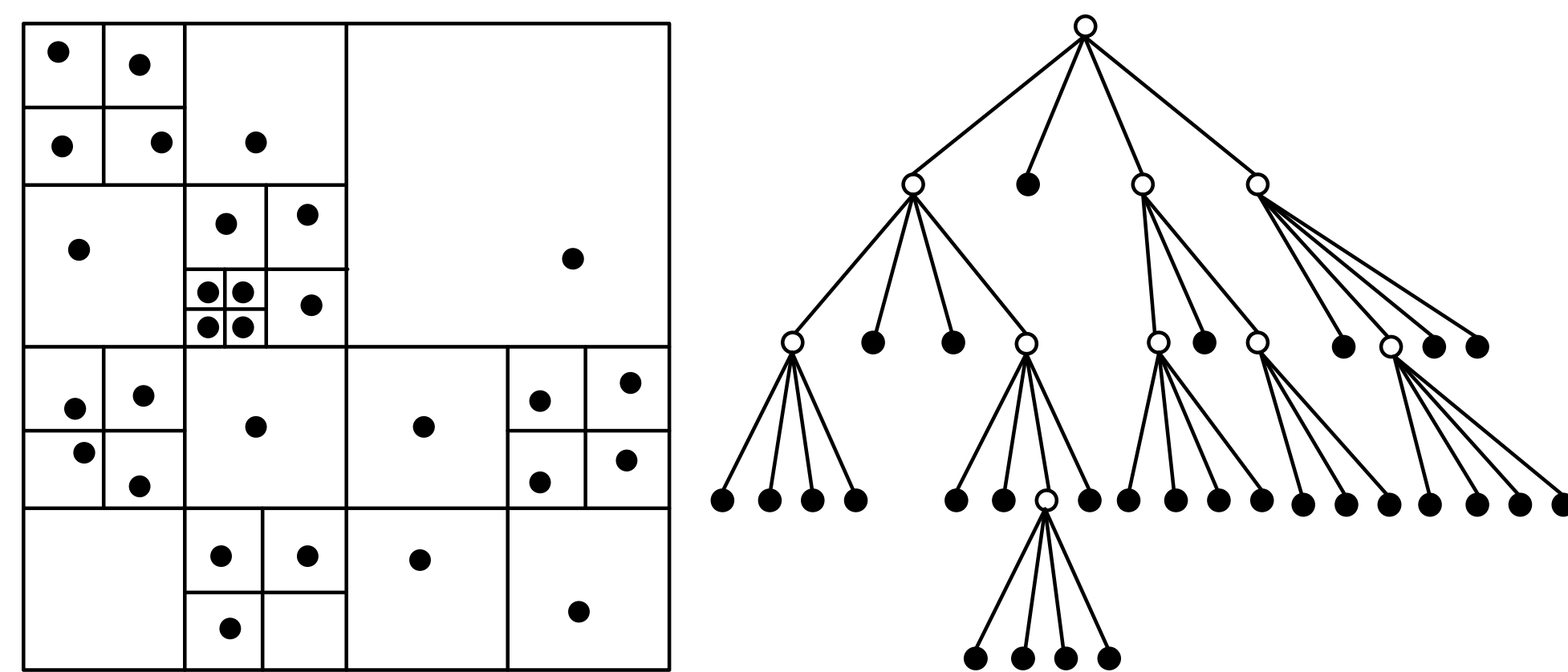


Fig. 1 A 2D body distribution and the corresponding quadtree. To compute forces on a body, we need to traverse the tree from the root. If a cell to be accessed is on remote node, we'd better cache it locally after fetching it, since we may reuse it for subsequent bodies.

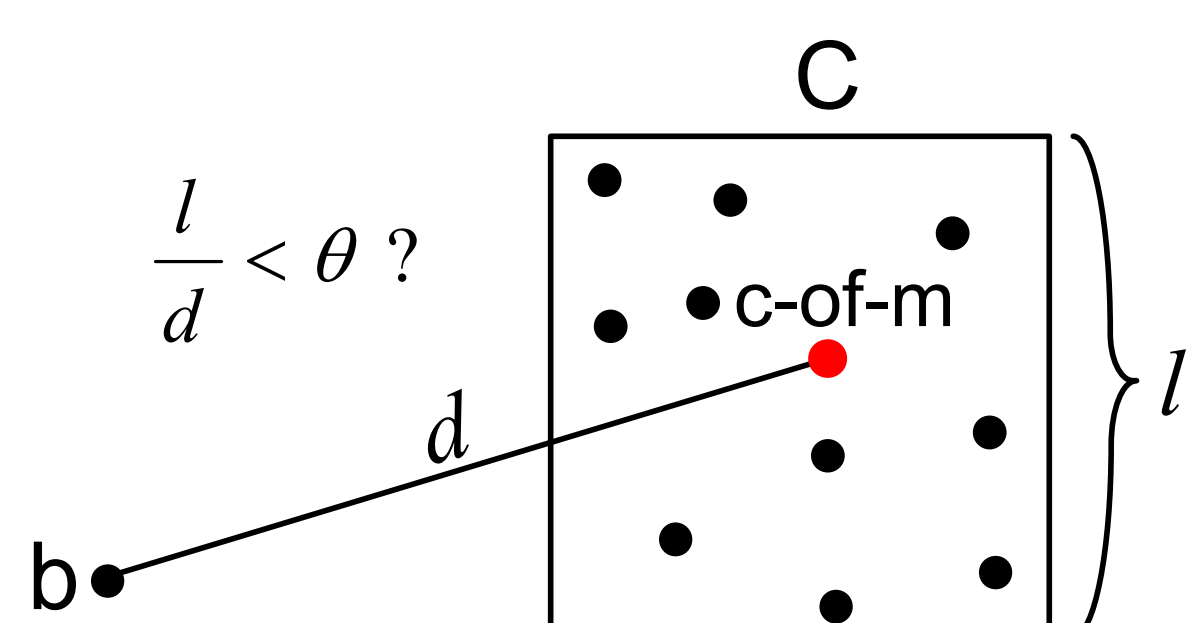


Fig. 2 Test if body b is far-away enough from cell C . If not, we need to "open" C . c-of-m = center-of-mass. θ is a user-defined constant. This test result is known at runtime, which makes it hard to reason communication statically.

Hybrid BH Design

- Implemented in PPL, a C++ template library having a PGAS memory model.
- Octree is distributed among compute nodes and cells are linked by *global* pointers provided by PPL.
- Each cell has a flag `localized` to indicate whether its children have been cached locally.

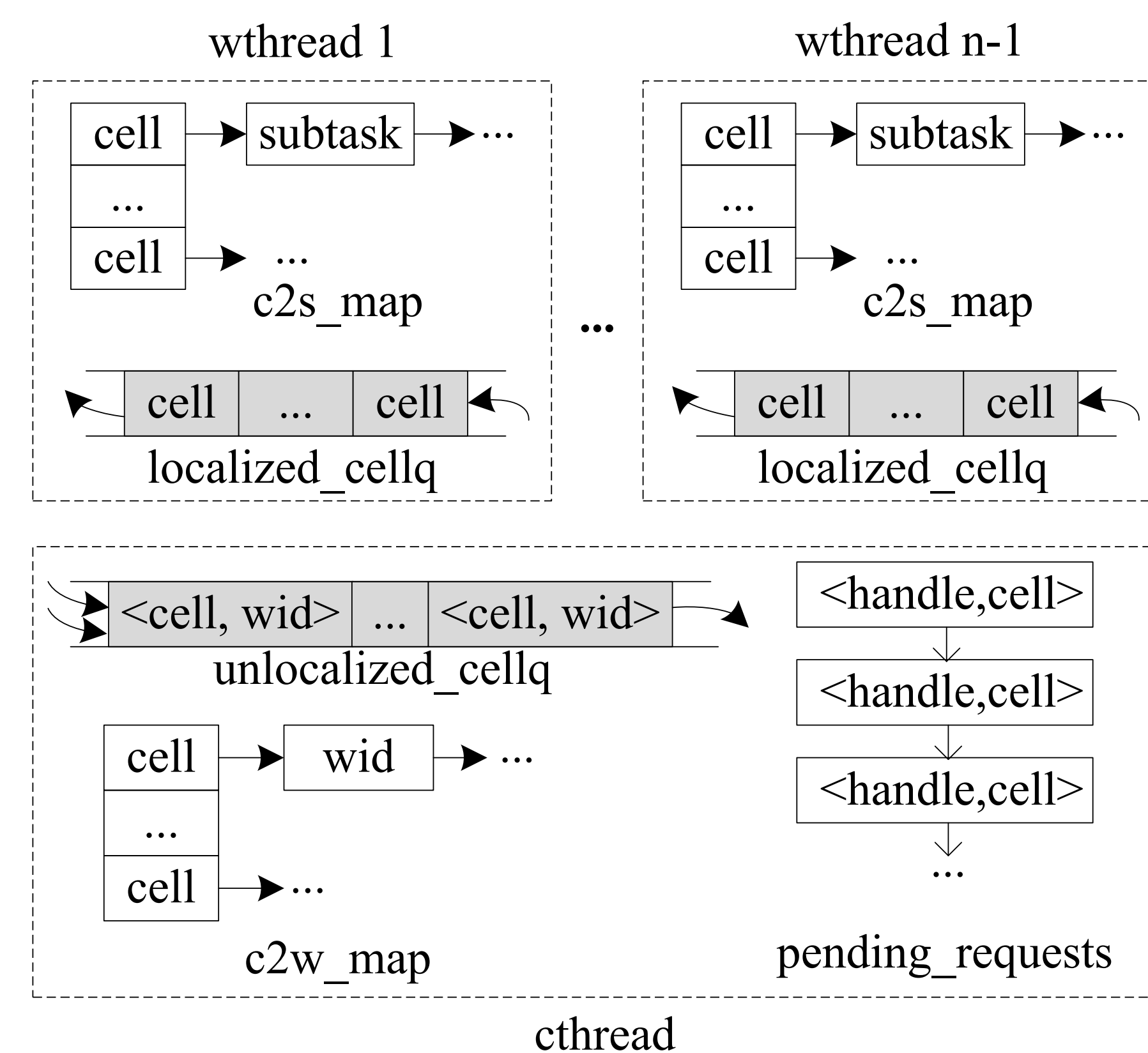


Fig. 3 Data structures used by threads, shaded are shared

- Multithreading:** One communication thread (cthread) + multiple worker threads (wthreads) per process.
- Multitasking** for fine-grain parallelism.
 - A task consists of tree traversal and force computation for one or more (if use tiling[4]) bodies;
 - A *subtask* encapsulates interactions between a body and a subtree rooted at an *unlocalized* cell.
- Wthreads** are responsible for computation.
 - Grab bodies from a body array assigned to the owner process and generate tasks.
 - Execute tasks (i.e., traverse the octree), whenever need to open a unlocalized cell, generate a subtask and continue to execute ready-to-run subtasks if any, or just generate a new task.
 - Manage subtasks locally through a cell-to-subtask map `c2s_map`.
 - Notify cthread which cells to localize by pushing `<cell, worker id>` to a queue `unlocalized_cellq`
- Cthread** is responsible for communication.
 - Pops cells from `unlocalized_cellq` and fetches their children by *dereferencing child pointers* (in a non-blocking fashion).
 - Caches child cells through *pointer-swizzling*.
 - Tests pending requests, if completed, sets corresponding cells as localized and pushes them back to `localized_cellq` on relevant wthreads
- Locality:** Between nodes, sort bodies in a space filling curve to reduce communications; Within a node, use tiling [4] to improve L1 cache reuse.

Experimental Setup

- Tested on a Linux cluster with each node having two hex-core Xeon® X5650 CPU@2.66GHz, with up to 64 nodes.
- Generated input bodies using the Plummer model.
- Compared the current code (named PPL BH) with UPC BH [1] (a BH code without multithreading), SPLASH-2 BH [5] (on single node) and other BH codes.

Results

- Less memory** by sharing the cached octree through multithreading instead of having each core build its own cached octree as in UPC BH.
 - With $n=1M$, $\theta=0.5$ on 64 nodes, saved 59% of memory.
- Less communication**, since data sharing also helps to filter duplicated off-node data requests.
 - With the same setting, saved 57% internode message.
- Better load balancing** by applying OpenMP *guided*-like scheduling, to adapt to runtime fluctuations.
 - Execution time variation of the 12 threads is $< 2.5\%$, which was up to 71% in UPC BH.
- Better performance and scalability**

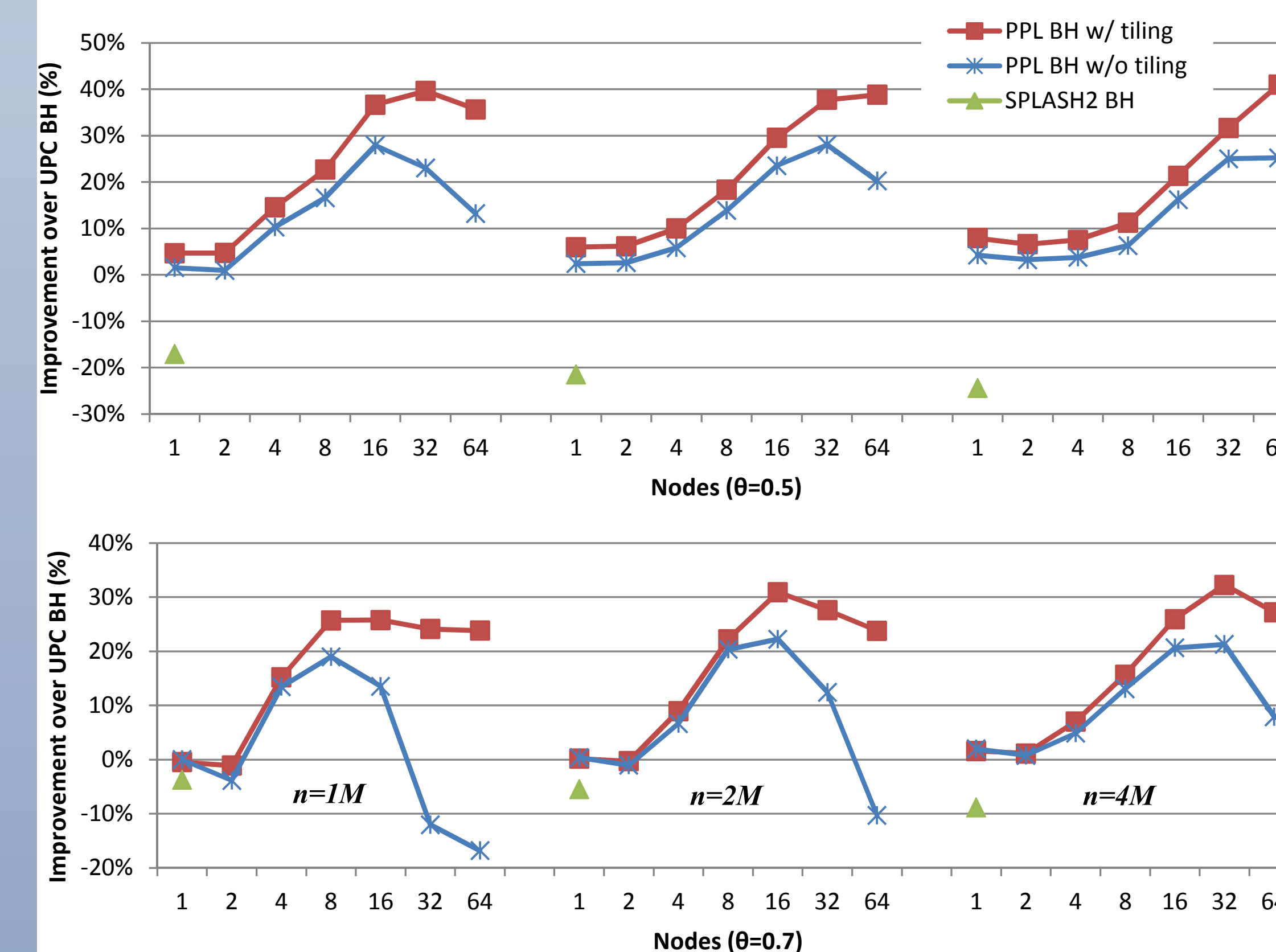


Fig. 4 Performance improvement over UPC BH. Note that the single node performance is even better than SPLASH-2 BH. With tiling, there're 128 bodies per task. Without tiling, 1 body per task.

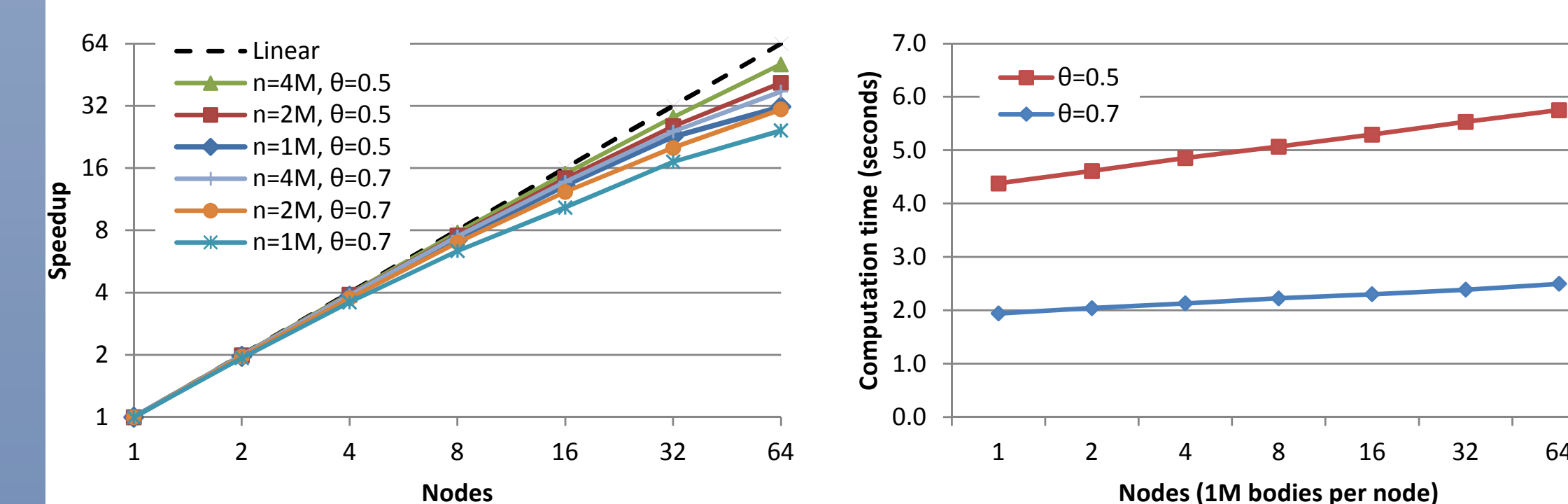


Fig. 5 Strong scaling (left) and weak scaling (right) curves.

(continue)

- Compared with two other state-of-the-art BH codes, i.e., Charm++ BH [3] and PEPC [2]
 - With $n=1M$, $\theta=0.5$, on 1~64 nodes, PPL BH is 1.4x ~ 4.5x faster, while its number of lines of code is 2x ~ 8x fewer.

Conclusions and Future Work

- We optimized the Barnes-Hut algorithm in a PGAS + X programming style, which brings many benefits as shown and thus is well suited to multicore clusters.
- We integrated multithreading and one-sided communication, with multitasking as the primary mechanism for latency hiding.
- In future, we like to abstract common runtime services from PPL BH and try to benefit other irregular applications.

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Acknowledgments

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